

# CIS 761. Database Management Systems

## Lecture notes on “The Chase”

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### A typical problem

Assume that the relation  $r$  satisfies the functional dependency

$$A \rightarrow B$$

and the multi-valued dependency

$$B \twoheadrightarrow C.$$

We want to use the above dependencies to *simplify* the query  $q$  given by

$$q = \Pi_{ABD}(r) \bowtie \Pi_{AC}(r).$$

### The solution

Assume that  $t$  is a tuple in  $q$ . Let  $t = (a, b, c, d)$ , then we know that  $(a, b, d) \in \Pi_{ABD}(r)$  and that  $(a, c) \in \Pi_{AC}(r)$ . Therefore, there exists  $c', b', d'$  such that  $r$  contains the tuples

$$\begin{array}{cccc} A & B & C & D \\ \hline a & b & c' & d \\ a & b' & c & d' \end{array}$$

Since  $r$  satisfies  $A \rightarrow B$ , we infer that  $b = b'$ , and the situation is therefore that  $r$  contains the tuples

$$\begin{array}{cccc} A & B & C & D \\ \hline a & b & c' & d \\ a & b & c & d' \end{array}$$

and since  $r$  satisfies  $B \twoheadrightarrow C$ ,  $r$  also contains the tuples

$A$	$B$	$C$	$D$
$a$	$b$	$c'$	$d$
$a$	$b$	$c$	$d'$
$a$	$b$	$c$	$d$
$a$	$b$	$c'$	$d'$

In particular, we see that  $t = (a, b, c, d) \in r$ . Since  $t$  was an arbitrary tuple in  $q$ , this shows that  $q \subseteq r$ . Clearly,  $r \subseteq q$ , so  $q = r$ . Thus, the complex query  $q$  can be reduced to the simple query  $r$ .

For more material on “the chase”, see [Abiteboul & Hull & Vianu, 1995].