



FROM INFORMATION TO MEANING

Recent Advances Concerning OWL and Rules

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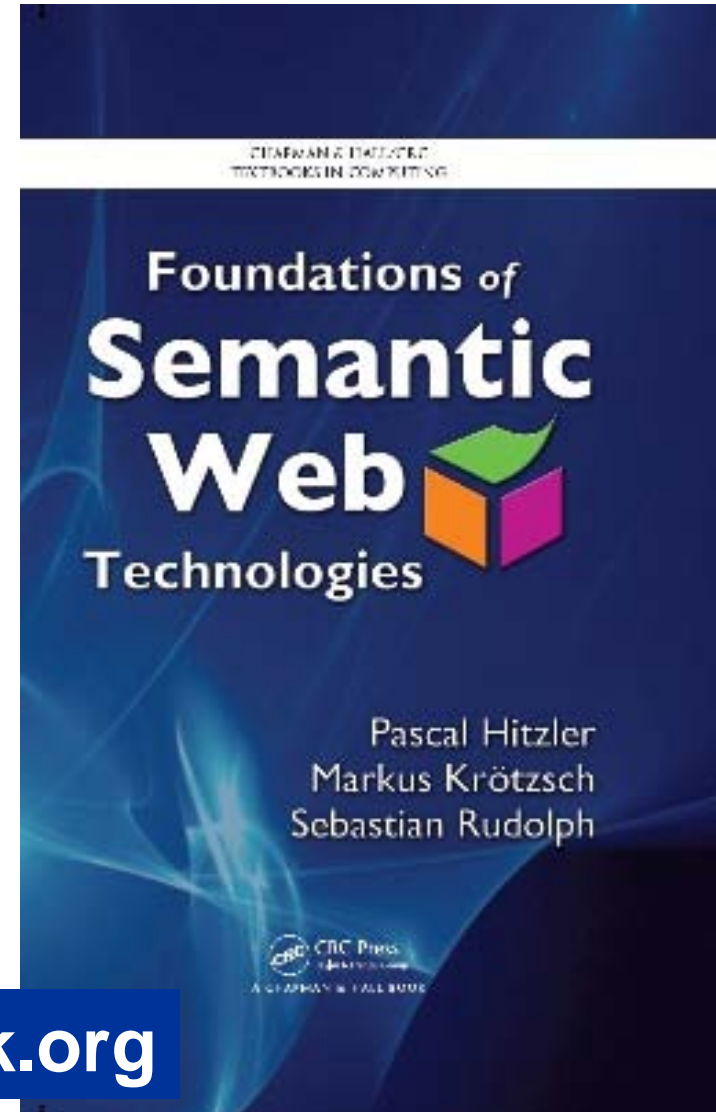
**Pascal Hitzler, Markus Krötzsch,
Sebastian Rudolph**

**Foundations of Semantic Web
Technologies**

Chapman & Hall/CRC, 2010

**Choice Magazine Outstanding Academic
Title 2010 (one out of seven in Information
& Computer Science)**

<http://www.semantic-web-book.org>



Pascal Hitzler, Markus Krötzsch, Sebastian Rudolph

语义Web技术基础

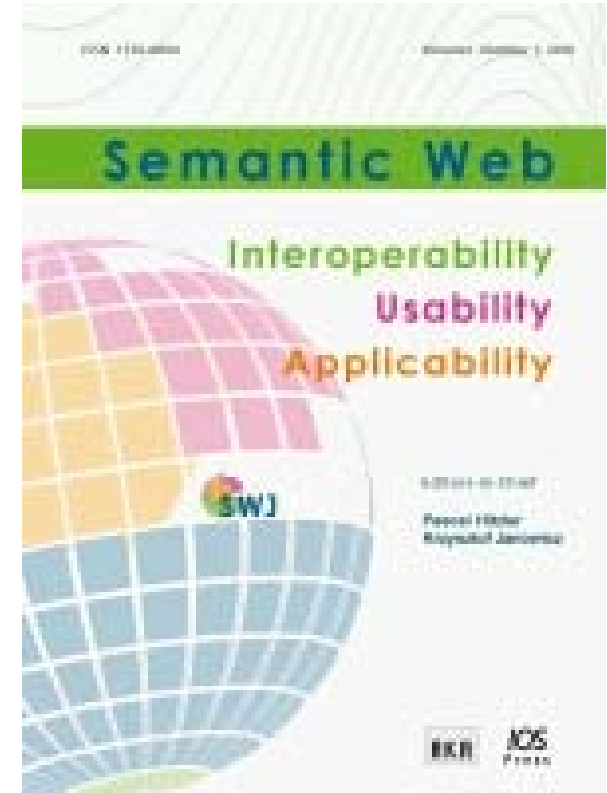
Tsinghua University Press (清华大学出版社), 2013.

Translators:

Yong Yu, Haofeng Wang, Guilin Qi (俞勇, 王昊奋, 漆桂林)

<http://www.semantic-web-book.org>

- **EiCs:** Pascal Hitzler
Krzysztof Janowicz
- **New journal with significant initial uptake.**
- **We very much welcome contributions at the “rim” of traditional Semantic Web research – e.g., work which is strongly inspired by a different field.**
- **Non-standard (open & transparent) review process.**



- **<http://www.semantic-web-journal.net/>**

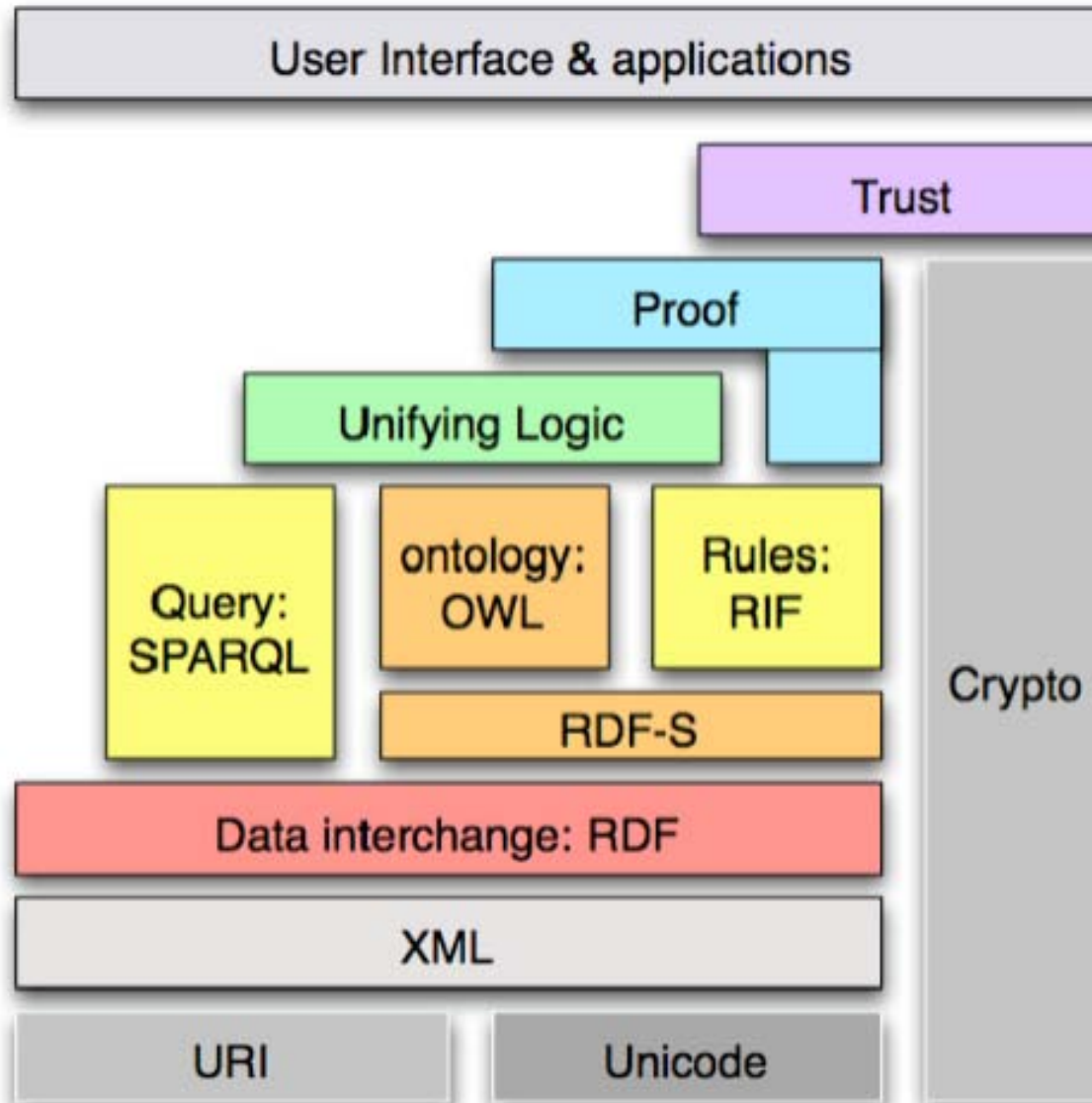
- **Ohio Center of Excellence in Knowledge-enabled Computing**
Director: Amit Sheth
- **15 faculty (8 in Computer Science) across 4 Departments, with ca. 50 PhD students**
- **Knowledge Engineering Lab (since January 2010)**
Led by myself

Currently **8 PhD students**
 2 Master students
 3 undergrads

- <http://www.knoesis.org/>



The Semantic Web Stack



- 1. Description Logics and OWL**
- 2. Rules expressible in description logics**
- 3. Extending description logics with rules through nominal schemas**
- 4. Algorithmizations for nominal schemas**
- 5. Adding non-monotonicity**
- 6. Conclusions**

Web Ontology Language (OWL)

- **W3C Recommendation since 2004**
- **OWL 2 since 2009**
- **based on description logics**
- **essentially, a decidable fragment of first-order predicate logic**

classes/concepts

A, B, C

unary predicates

A(x), B(x), C(x)

roles/properties

R, S

binary predicates

R(x,y), S(x,y)

individuals

a, b, c

constants

a, b, c

class conjunction

$$C \sqcap D$$

$$C(x) \wedge D(x)$$

existential restriction

$$\exists R.C$$

$$\exists y (R(x,y) \wedge C(y))$$

class inclusion/subsumption

$$C \sqsubseteq D$$

$$C(x) \rightarrow D(x)$$

$$C \equiv D$$

$$C(x) \leftrightarrow D(x)$$

role chains

$$R_1 \circ \dots \circ R_n \sqsubseteq R$$

$$R_1(x, x_1) \wedge \dots \wedge R(x_n, x_{n+1}) \rightarrow R(x, x_{n+1})$$

ThaiDish $\sqsubseteq \exists$ contains.Nut

Nutallergic $\sqcap \exists$ eats.Nut \sqsubseteq Unhappy

eats \circ contains \sqsubseteq eats

inverse roles

$$R \equiv S^{-}$$

$$R(x,y) \leftrightarrow S(y,x)$$

This logic is already undecidable!

(see e.g. [ISWC 2007])

Name of the logic: ELRI

Decidability is a central characteristics of description logics.

- 1. Disallow \exists :**
Essentially leads to OWL RL.
Fragment of Datalog.
Tractable (i.e., polynomial complexity).
- 2. Disallow inverse roles:**
Essentially leads to OWL EL.
Akin “in spirit” to existential rules/Datalog+-.
Tractable.
- 3. Restrict recursion in role chains (a.k.a. *regularity* restriction):**
With further constructors, leads to OWL DL, a.k.a. SROIQ.
Decidable, but not tractable.

The following can be used in OWL EL (logic remains tractable).

Self

$C \sqsubseteq \exists R.\text{Self}$ $C(x) \rightarrow R(x,x)$

Can be used e.g. for typecasting.

nominals

$\{a\} \sqsubseteq C$

$C(a)$

a is a constant

$C \sqsubseteq \{a\}$

$C(x) \rightarrow x=a$

$\{a\} \equiv \{b\}$

$\rightarrow a=b$

$A \sqcap \exists R.\{b\} \sqsubseteq C$ becomes $A(x) \wedge R(x,b) \rightarrow C(x)$

The following are used in expressive (intractable) DLs

class negation

$\neg C$

$\neg C(x)$

class disjunction

$C \sqcup D$

$C(x) \vee D(x)$

universal restriction

$\forall R.C$

$\forall y (R(x,y) \rightarrow C(y))$

There are some more of course.

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Which rules can be encoded in OWL?

$A \sqsubseteq B$ becomes $A(x) \rightarrow B(x)$

$R \sqsubseteq S$ becomes $R(x, y) \rightarrow S(x, y)$

$A \sqcap \exists R. \exists S. B \sqsubseteq C$ becomes $A(x) \wedge R(x, y) \wedge S(y, z) \wedge B(z) \rightarrow C(x)$

$\{a\} \equiv \{b\}$ becomes $\rightarrow a = b.$

$A \sqcap B \sqsubseteq \perp$ becomes $A(x) \wedge B(x) \rightarrow f.$

$R \circ S \sqsubseteq T$ becomes $R(x, y) \wedge S(y, z) \rightarrow T(x, z)$

Which rules can be encoded in OWL?

$A \sqsubseteq \neg B \sqcup C$ becomes $A(x) \wedge B(x) \rightarrow C(x)$

$A \sqsubseteq \forall R.B$ becomes $A(x) \wedge R(x, y) \rightarrow B(y)$

$A \sqsubseteq B \wedge C$ becomes $A(x) \rightarrow B(x)$ and $A(x) \rightarrow C(x)$

$A \sqcup B \rightarrow C$ becomes $A(x) \rightarrow C(x)$ and $B(x) \rightarrow C(x)$

$$\text{Elephant}(x) \wedge \text{Mouse}(y) \rightarrow \text{biggerThan}(x, y)$$

- **Rolification of a concept A:** $A \equiv \exists R_A.\text{Self}$

$$\text{Elephant} \equiv \exists R_{\text{Elephant}}.\text{Self}$$

$$\text{Mouse} \equiv \exists R_{\text{Mouse}}.\text{Self}$$

$$R_{\text{Elephant}} \circ U \circ R_{\text{Mouse}} \sqsubseteq \text{biggerThan}$$

$A(x) \wedge R(x, y) \rightarrow S(x, y)$ becomes $R_A \circ R \sqsubseteq S$

$A(y) \wedge R(x, y) \rightarrow S(x, y)$ becomes $R \circ R_A \sqsubseteq S$

$A(x) \wedge B(y) \wedge R(x, y) \rightarrow S(x, y)$ becomes $R_A \circ R \circ R_B \sqsubseteq S$

$\text{Woman}(x) \wedge \text{marriedTo}(x, y) \wedge \text{Man}(y) \rightarrow \text{hasHusband}(x, y)$

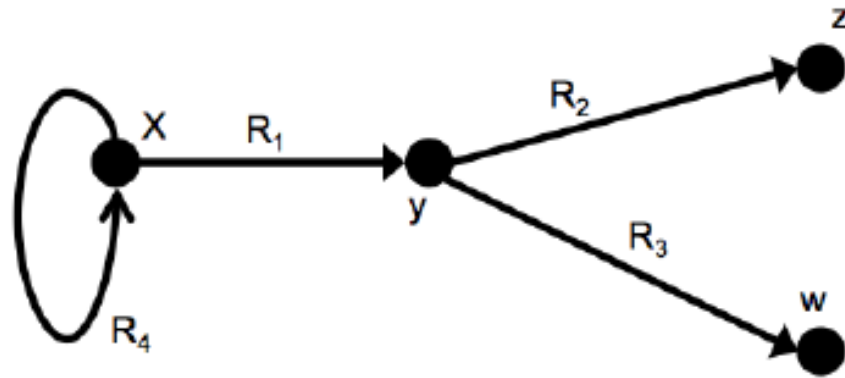
$R_{\text{Woman}} \circ \text{marriedTo} \circ R_{\text{Man}} \sqsubseteq \text{hasHusband}$

careful – regularity of RBox needs to be retained:

$\text{hasHusband} \sqsubseteq \text{marriedTo}$

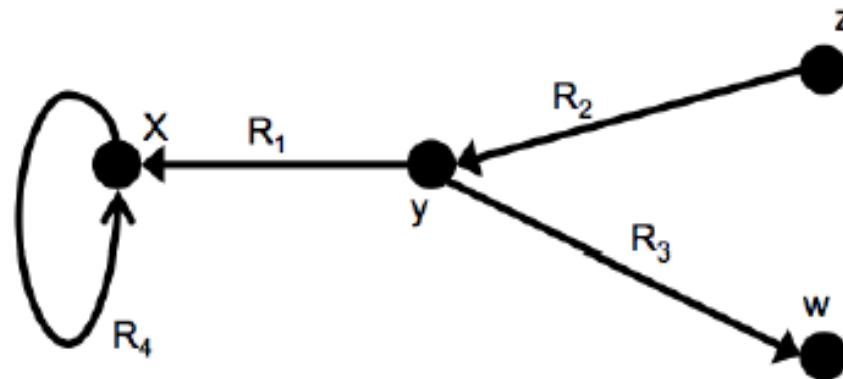
$$\text{worksAt}(x, y) \wedge \text{University}(y) \wedge \text{supervises}(x, z) \wedge \text{PhDStudent}(z) \\ \rightarrow \text{professorOf}(x, z)$$
$$R_{\exists \text{worksAt.University}} \circ \text{supervises} \circ R_{\text{PhDStudent}} \sqsubseteq \text{professorOf.}$$

$$R_1(x, y) \wedge C_1(y) \wedge R_2(y, w) \wedge R_3(y, z) \wedge C_2(z) \wedge R_4(x, x) \rightarrow C_3(x)$$



$$\exists R_1. (C_1 \sqcap \exists R_2. \top \sqcap \exists R_3. C_2) \sqcap \exists R_4. \text{Self} \sqsubseteq C_3$$

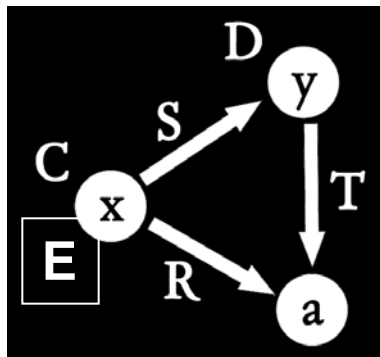
$$R_1(y, x) \wedge C_1(y) \wedge R_2(w, y) \wedge R_3(y, z) \wedge C_2(z) \wedge R_4(x, x) \rightarrow C_3(x)$$



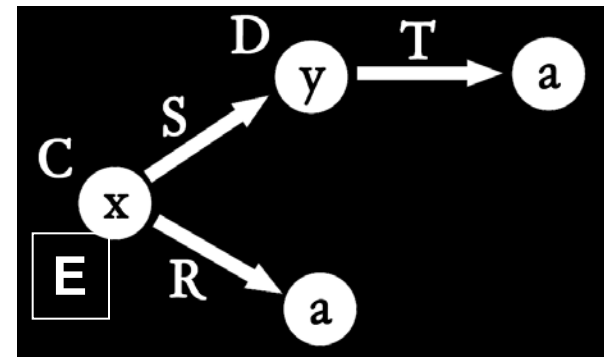
$$\exists R_1^-. (C_1 \sqcap \exists R_2^-. \top \sqcap \exists R_3. C_2) \sqcap \exists R_4. Self \sqsubseteq C_3$$

So how can we pinpoint this?

- Tree-shaped bodies
- First argument of the conclusion is the root
- $C(x) \wedge R(x,a) \wedge S(x,y) \wedge D(y) \wedge T(y,a) \rightarrow E(x)$
 - $C \sqcap \exists R.\{a\} \sqcap \exists S.(D \sqcap \exists T.\{a\}) \sqsubseteq E$



duplicating
nominals
is
ok



$$C(x) \wedge R(x, a) \wedge S(x, y) \wedge D(y) \wedge T(y, a) \rightarrow P(x, y)$$

$$a_1 \longleftarrow x \longrightarrow y \longrightarrow a_2$$

$$C \sqcap \exists R.\{a\} \sqsubseteq \exists R1.\text{Self}$$

$$D \sqcap \exists T.\{a\} \sqsubseteq \exists R2.\text{Self}$$

$$R1 \circ S \circ R2 \sqsubseteq P$$

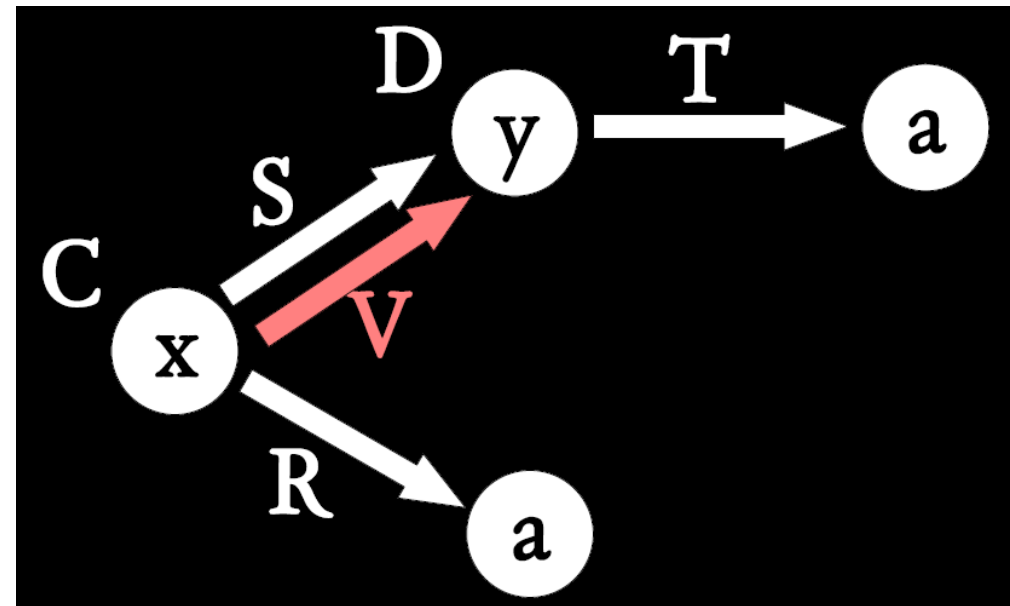
So how can we pinpoint this?

- Tree-shaped bodies
- First argument of the conclusion is the root
- $C(x) \wedge R(x,a) \wedge S(x,y) \wedge D(y) \wedge T(y,a) \rightarrow V(x,y)$

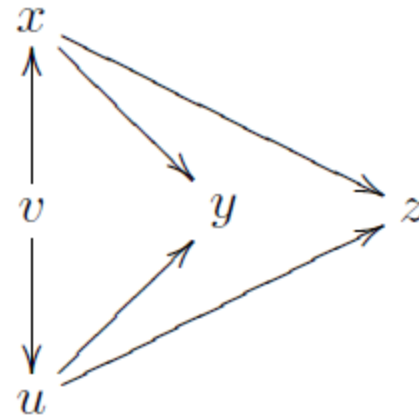
$C \sqcap \exists R.\{a\} \sqsubseteq \exists R1.Self$

$D \sqcap \exists T.\{a\} \sqsubseteq \exists R2.Self$

$R1 \circ S \circ R2 \sqsubseteq V$



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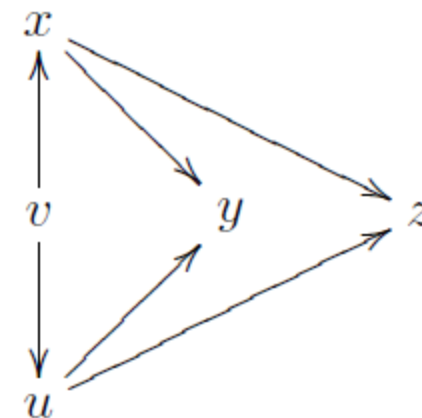
$$\begin{aligned} & \text{hasReviewAssignment}(v, x) \wedge \text{hasAuthor}(x, y) \wedge \text{atVenue}(x, z) \\ & \wedge \text{hasSubmittedPaper}(v, u) \wedge \text{hasAuthor}(u, y) \wedge \text{atVenue}(u, z) \\ & \rightarrow \text{hasConflictingAssignedPaper}(v, x) \end{aligned}$$


with **y,z constants:**

$$\begin{aligned} R_{\exists \text{hasSubmittedPaper}.(\exists \text{hasAuthor}.\{y\} \sqcap \exists \text{atVenue}.\{z\})} & \circ \text{hasReviewAssignment} \\ & \circ R_{\exists \text{hasAuthor}.\{y\} \sqcap \exists \text{atVenue}.\{z\}} \\ & \sqsubseteq \text{hasConflictingAssignedPaper} \end{aligned}$$

$$\begin{aligned} & \text{hasReviewAssignment}(v, x) \wedge \text{hasAuthor}(x, y) \wedge \text{atVenue}(x, z) \\ & \wedge \text{hasSubmittedPaper}(v, u) \wedge \text{hasAuthor}(u, y) \wedge \text{atVenue}(u, z) \\ & \rightarrow \text{hasConflictingAssignedPaper}(v, x) \end{aligned}$$

assume y, z bind only to named individuals
we introduce a new construct, called
nominal schemas
or *nominal variables*



$R_{\exists \text{hasSubmittedPaper}.(\exists \text{hasAuthor}.\{y\} \sqcap \exists \text{atVenue}.\{z\})}$ \circ hasReviewAssignment
 $\circ R_{\exists \text{hasAuthor}.\{y\} \sqcap \exists \text{atVenue}.\{z\}}$
 \sqsubseteq hasConflictingAssignedPaper

$$\text{hasChild}(x, y) \wedge \text{hasChild}(x, z) \wedge \text{classmate}(y, z) \rightarrow C(x)$$
$$\exists \text{hasChild}.\{z\} \sqcap \exists \text{hasChild}.\exists \text{classmate}.\{z\} \sqsubseteq C$$

- Decidability is retained.
- Complexity is *the same*.

- A naïve implementation is straightforward:

Replace every axiom with nominal schemas by a set of OWL 2 axioms, obtained from *grounding* the nominal schemas.

However, this may result in a lot of new OWL 2 axioms.
The naïve approach will probably only work for ontologies with *few* nominal schemas.

- A powerful macro.
- A conceptual bridge to rule formalism:

We can actually also express all DL-safe Datalog rules!

$$R(x, y) \wedge A(y) \wedge S(z, y) \wedge T(x, z) \rightarrow P(z, x)$$

$$\begin{aligned} & \exists U.(\{x\} \sqcap \exists R.\{y\}) \\ & \sqcap \exists U.(\{y\} \sqcap A) \\ & \sqcap \exists U.(\{z\} \sqcap \exists S.\{y\}) \\ & \sqcap \exists U.(\{x\} \sqcap \exists T.\{z\}) \\ & \sqsubseteq \exists U.(\{z\} \sqcap \exists P.\{x\}) \end{aligned}$$

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Naïve implementation – experiments

	No axioms added		1 different ns		2 different ns		3 different ns	
Fam (5)	0.01''	0.00''	0.01''	0.00''	0.01''	0.00''	0.04''	0.02''
Swe (22)	3.58''	0.08''	3.73''	0.07''	3.85''	0.10''	10.86''	1.11''
Bui (42)	2.7''	0.16''	2.5''	0.15''	2.75''	0.26''	1' 14'	6.68''
Wor (80)	0.11''	0.04''	0.12''	0.05''	1.1''	0.55''	OOM *	OOM*
Tra (183)	0.05''	0.03''	0.05''	0.02''	5.66''	1.76''	OOM	OOM
FTr (368)	0.03''	4.28''	0.05	5.32''	35.53''	42.73''	OOM	OOM
Eco (482)	0.04''	0.24''	0.07''	0.02''	56.59''	13.67''	OOM	OOM

OOM = Out of Memory

from the TONES repository:

Ontology	Classes	Data P.	Object P.	Individuals
Fam	4	1	11	5
Swe	189	6	25	22
Bui	686	0	24	42
Wor	1842	0	31	80
Tra	445	4	89	183
FTr	22	6	52	368
Eco	339	8	45	482

- **Adding nominal schemas to existing tableaux algorithms:**

grounding : if $C \in L(s)$, $\{z\}$ is a nominal schema in C ,
 $C[z/a_i] \notin L(s)$ for some $i, 1 \leq i \leq \ell$
then $L(s) := L(s) \cup \{C[z/a_i]\}$

plus some restrictions on existing tableaux rules, essentially to ensure that (1) no variable binding is broken and (2) nominal schemas are not propagated through the tableau.

- variant of absorption [Steigmiller, Glimm, Liebig, IJCAI-13]
- essentially, a sort of smart rewriting as pre-processing

Example 1 *Our running example $\exists r.(\{x\} \sqcap \exists a.\{y\} \sqcap \exists v.\{z\}) \sqcap \exists s.(\exists a.\{y\} \sqcap \exists v.\{z\}) \sqsubseteq \exists c.\{x\}$ can be almost completely absorbed into the following axioms:*

$$\begin{array}{lll} O \sqsubseteq \downarrow x.T_x & T_z \sqsubseteq \forall v^-.T_2 & (T_1 \sqcap T_2) \sqsubseteq T_3 \\ O \sqsubseteq \downarrow y.T_y & T_3 \sqsubseteq \forall s^-.T_4 & (T_3 \sqcap T_x) \sqsubseteq T_5 \\ O \sqsubseteq \downarrow z.T_z & T_5 \sqsubseteq \forall r^-.T_6 & (T_4 \sqcap T_6) \sqsubseteq T_7. \\ T_y \sqsubseteq \forall a^-.T_1 & T_7 \sqsubseteq gr(\exists c.\{x\}), & \end{array}$$

where $T_x, T_y, T_z, T_1, \dots, T_7$ are fresh atomic concepts. Only $\exists c.\{x\}$ cannot be absorbed and has to be grounded on demand.

[Steigmiller, Glimm, Liebig, IJCAI-13]

Table 2: DL-safe Rules for UOBM-Benchmarks

Name	DL-safe Rule	Matches
R1	$isFirendOf(?x, ?y), like(?x, ?z), like(?y, ?z) \rightarrow hasLink1(?x, ?y)$	4,037
R2	$isFirendOf(?x, ?y), takesCourse(?x, ?z), takesCourse(?y, ?z) \rightarrow hasLink2(?x, ?y)$	82
R3	$takesCourse(?x, ?z), takesCourse(?y, ?z), hasSameHomeTownWith(?x, ?y) \rightarrow hasLink3(?x, ?y)$	940
R4	$hasDoctoralDegreeFrom(?x, ?z), hasMasterDegreeFrom(?x, ?w), hasDoctoralDegreeFrom(?y, ?z), hasMasterDegreeFrom(?y, ?w), worksFor(?x, ?v), worksFor(?y, ?v), \rightarrow hasLink4(?x, ?y)$	369
R5	$isAdvisedBy(?x, ?z), isAdvisedBy(?y, ?z), like(?x, ?w), like(?y, ?w), like(?z, ?w) \rightarrow hasLink5(?x, ?y)$	286

Table 3: Comparison of the increases in reasoning time of the consistency tests for $UOBM_1 \setminus D$ extended by rules in seconds

Rule	upfront grounding		direct propagation		representative propagation		Hermit	Pellet
		mem	without BC	with BC	without BC	with BC		
R1	(10.99)	mem	9.12	7.10	5.06	3.38	31.46	6.33
R2	(10.92)	4.05	3.33	2.33	2.13	2.11	4.79	7.4
R3	(13.33)	3.55	1.98	0.62	2.20	0.76	1.67	142.25
R4	(16.44)	0.30	1.08	0.09	1.06	0.07	1.42	122.85
R5	(time)	–	1.87	0.50	1.80	0.43	28.41	mem

Based on [Krötzsch, JELIA10]

Ontology	Individuals	no ns	1 ns	2 ns	3 ns	4 ns	5 ns
Rex (full ground.)	100	263	263 (321)	267 (972)	273	275	259
	1000	480	518 (1753)	537 (OOM)	538	545	552
	10000	2904	2901 (133179)	3120 (OOM)	3165	3192	3296
Spatial (full ground.)	100	22	191 (222)	201 (1163)	198	202	207
	1000	134	417 (1392)	415 (OOM)	421	431	432
	10000	1322	1792 (96437)	1817 (OOM)	1915	1888	1997
Xenopus (full ground.)	100	62	332 (383)	284 (1629)	311	288	280
	1000	193	538 (4751)	440 (OOM)	430	456	475
	10000	1771	2119 (319013)	1843 (OOM)	1886	2038	2102

- We rewrite mincardinality restrictions into maxcardinality restrictions or approximate using an existential.
- We rewrite universal quantification into existential quantification.
- We approximate maxcardinality restrictions using functionality.
- We approximate inverse roles and functionality using nominal schemas.
- We approximate negation using class disjointness.
- We approximate disjunction using conjunction.

- **inverses:** $\{x\} \sqcap \exists R.\{y\} \sqsubseteq \{y\} \sqcap \exists S.\{x\}$

- **functionality** $C \sqsubseteq \leq 1 R.D$:

$$C \sqcap \exists R.(\{z1\} \sqcap D) \sqcap \exists R.(\{z2\} \sqcap D) \sqsubseteq \exists U.(\{z1\} \sqcap \{z2\})$$

Approximation results (using IRIS)

Ontology	HermiT	Fact++	Pellet	Ours	Ours Recall
BAMS	3	2	10	107	100%
DOLCE	1	1	4	53	100%
GALEN	4	2	17	7840	90.8%
GO	36	75	59	N/A	N/A
GardinerCorpus	14	6	17	89	92.3%
OBO	34	61	139	N/A	N/A

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- **[Knorr, Hitzler, Maier ECAI2012]**
- **Extension of an autoepistemic description logic approach by nominal schemas.**
- **Results in a language which incorporates most of the major approaches to non-monotonic extensions of DLs.**
- **E.g. covers**
 - hybrid MKNF [Motik & Rosati], which in turn covers
 - non-disjunctive ASP
 - DL Programs / dlvhex (Eiter et al.)
- **Also covers OWL / SROIQ(D) of course.**

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- **Paradigms are converging.**
- **More work needed e.g. re.**
 - **algorithmizations**
 - **relating OWL EL and existential rules research**
 - **making non-monotonic reasoning fit for semantic web applications**



Collaborators on the covered topics

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A tutorial:

- **Adila A. Krisnadhi, Frederick Maier, Pascal Hitzler, OWL and Rules. In: A. Polleres, C. d'Amato, M. Arenas, S. Handschuh, P. Kroner, S. Ossowski, P.F. Patel-Schneider (eds.), Reasoning Web. Semantic Technologies for the Web of Data. 7th International Summer School 2011, Galway, Ireland, August 23-27, 2011, Tutorial Lectures. Lecture Notes in Computer Science Vol. 6848, Springer, Heidelberg, 2011, pp. 382-415.**

Background reading:

- **Pascal Hitzler, Markus Krötzsch, Sebastian Rudolph, Foundations of Semantic Web Technologies. Textbooks in Computing, Chapman and Hall/CRC Press, 2009.
<http://www.semantic-web-book.org/>**
- **Pascal Hitzler, Markus Krötzsch, Bijan Parsia, Peter F. Patel-Schneider, Sebastian Rudolph, OWL 2 Web Ontology Language: Primer (Second Edition). W3C Recommendation, 11 December 2012. <http://www.w3.org/TR/owl2-primer/>**

- **Markus Krötzsch, Frederick Maier, Adila Alfa Krisnadhi, Pascal Hitzler, A Better Uncle For OWL – Nominal Schemas for Integrating Rules and Ontologies. In: S. Sadagopan, Krithi Ramamritham, Arun Kumar, M.P. Ravindra, Elisa Bertino, Ravi Kumar (eds.), WWW '11 20th International World Wide Web Conference, Hyderabad, India, March/April 2011. ACM, New York, 2011, pp. 645-654.**
- **Markus Krötzsch, Sebastian Rudolph, Pascal Hitzler, Description Logic Rules. In: Malik Ghallab, Constantine D. Spyropoulos, Nikos Fakotakis, Nikos Avouris (eds.), Proceedings of the 18th European Conference on Artificial Intelligence, ECAI2008, Patras, Greece, July 2008. IOS Press, 2008, pp. 80-84.**
- **Markus Krötzsch. Description Logic Rules. Studies on the Semantic Web, Vol. 008, IOS Press, 2010.
<http://www.semantic-web-studies.net/>**

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